

Commuting upper triangular binary morphisms

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Abstract

A morphism g from the free monoid X^* into itself is called upper triangular if the matrix of g is upper triangular. We characterize all upper triangular binary morphisms g_1 and g_2 such that $g_1g_2 = g_2g_1$.

Keywords: Free monoid morphism; Commutativity; Combinatorics on morphisms

1 Introduction

The free monoid morphisms play an important role in many areas of mathematics and theoretical computer science (see [1, 8, 9, 11, 12, 13]). On the other hand, many questions concerning combinatorics on morphisms appear to be rather difficult. It is instructive to consider the problem of commutativity. If u and v are words, the equation $uv = vu$ holds if and only if there is a word w such that u and v are powers of w (see [8]). For free monoid morphisms the situation is more complicated. For two morphisms g_1 and g_2 , the equation $g_1g_2 = g_2g_1$ does not imply that g_1 and g_2 are powers of a third morphism (see, however, [10]).

In this paper we study commuting upper triangular binary morphisms. Let $X = \{a, b\}$ be a binary alphabet. A morphism g from the free monoid X^* into itself is called upper triangular if the matrix of g is upper triangular. If a is the first letter of X , this means that there is a nonnegative integer s such that $g(a) = a^s$. We will characterize all upper triangular binary morphisms g_1 and g_2 such that $g_1g_2 = g_2g_1$.

We now outline the contents of this paper. In Section 2 we recall the basic definitions. In Section 3 we discuss the connections between freeness and commutativity. In Section 4 we give examples of commuting morphisms. In Section 5 we study infinite words generated by morphisms. While the morphisms we study are not uniform, it turns out to be possible to use results concerning automatic sequences. In particular, we will apply the theorem of Cobham character-

izing those sequences which are automatic in two multiplicatively independent bases (see [1]).

In Sections 6,7,8 and 9 we characterize all upper triangular binary morphisms g_1 and g_2 such that $g_1g_2 = g_2g_1$. Assume that a is the first letter and b is the second letter of the binary alphabet. In Section 6 we consider nonsingular morphisms such that both $g_1(b)$ and $g_2(b)$ contain at least two occurrences of b . We have two cases depending on whether these numbers are multiplicatively independent or not. The remaining cases are easier and are discussed in Sections 7,8 and 9.

We assume that the reader is familiar with the basics of free monoid morphisms, infinite words, automatic sequences and combinatorics on words (see [1, 8, 9, 11, 12, 13]). For previous results concerning combinatorics on morphisms see, e.g., [4, 5, 6, 7, 10].

2 Definitions

We use standard notation and terminology concerning free monoids and their morphisms (see [1, 8, 9, 11, 12]). If X is a finite nonempty set, X^* is the *free monoid* generated by X . The identity element of X^* is the *empty word* denoted by ε . If u, v, w are words such that $uv = w$, we denote $v = u^{-1}w$.

If w is a word and a is a letter, then $|w|_a$ is the number of occurrences of a in w . The *length* of a word w , denoted by $|w|$, is the total number of letters in w .

Let X and Y be finite nonempty alphabets. A mapping $h : X^* \rightarrow Y^*$ is a *morphism* if

$$h(uv) = h(u)h(v)$$

for all $u, v \in X^*$. The set of all morphisms from X^* to X^* is denoted by $\text{Hom}(X^*)$. $\text{Hom}(X^*)$ is a monoid with respect to the usual product of morphisms.

If $h \in \text{Hom}(X^*)$ and the letters of X are x_1, \dots, x_d in a fixed order, then the *matrix* M_h of h is defined by

$$M_h = \begin{pmatrix} |h(x_1)|_{x_1} & |h(x_2)|_{x_1} & \dots & |h(x_d)|_{x_1} \\ |h(x_1)|_{x_2} & |h(x_2)|_{x_2} & \dots & |h(x_d)|_{x_2} \\ \vdots & \vdots & & \vdots \\ |h(x_1)|_{x_d} & |h(x_2)|_{x_d} & \dots & |h(x_d)|_{x_d} \end{pmatrix}.$$

A morphism $h \in \text{Hom}(X^*)$ is *upper triangular* if its matrix M_h is upper triangular. The set of upper triangular morphisms from X^* to X^* is denoted by $\text{Tri}(X^*)$. A morphism $h \in \text{Hom}(X^*)$ is *nonsingular* if its matrix is nonsingular.

Let now X be a finite alphabet and let $h \in \text{Hom}(X^*)$. If $w \in X^*$ is a word such that w is a prefix of $h(w)$ and $\lim_{n \rightarrow \infty} |h^n(w)| = \infty$, we say that h is *prolongable* on w and define the infinite word $h^\omega(w)$ by

$$h^\omega(w) = \lim_{n \rightarrow \infty} h^n(w).$$

Hence, $h^\omega(w)$ is the unique infinite word u such that $h^n(w)$ is a prefix of u for all $n \in \mathbb{N}$.

3 Connections between freeness and commutativity

A nonempty subset Y of a semigroup S is called *free* if every element of the subsemigroup of S generated by Y can be written uniquely as a product of elements of Y . In other words, a set Y is free if for all positive integers m and n and $u_1, \dots, u_m, v_1, \dots, v_n \in Y$, the equation

$$u_1 u_2 \cdots u_m = v_1 v_2 \cdots v_n$$

implies that

$$m = n \quad \text{and} \quad u_i = v_i \quad \text{for} \quad i = 1, \dots, m.$$

For an excellent introduction to freeness problems over semigroups we refer to [2].

If a set contains two elements which commute, then the set is not free. If $u, v \in X^*$ and $u \neq v$, then $\{u, v\}$ is free if and only if u and v do not commute (see [8]).

We recall some related results for upper triangular morphisms.

First, let $X = \{a, b\}$ be a binary alphabet. Let $g_1, g_2 \in \text{Tri}(X^*)$. We say that $\{g_1, g_2\}$ is a *special pair* if $g_1(b)$ and $g_2(b)$ belong to a^*ba^* and exactly one of $g_1(a)$ and $g_2(a)$ equals a .

The following result is from [7].

Theorem 1 *Let $X = \{a, b\}$ and let $g_1, g_2 \in \text{Tri}(X^*)$ be nonsingular upper triangular morphisms. Assume that $g_1 \neq g_2$. Assume that $\{g_1, g_2\}$ is not a special pair. If $\{g_1, g_2\}$ is not free, then $g_1 g_2 = g_2 g_1$.*

For larger alphabets we have the following result (see [6]).

Theorem 2 *Let X be an arbitrary alphabet. Let $g_1, g_2 \in \text{Tri}(X^*)$ and let M_i be the matrix of g_i for $i = 1, 2$. Assume $g_1 \neq g_2$. Assume that all diagonal entries of M_i are at least two for $i = 1, 2$. If $\{g_1, g_2\}$ is not free, then $g_1 g_2 = g_2 g_1$.*

Theorems 1 and 2 imply the following lemma.

Lemma 3 *Assume that the morphisms g_1 and g_2 satisfy the assumptions of Theorem 1 or Theorem 2. Let m and n be positive integers. If g_1^m and g_2^n commute, then g_1 and g_2 commute.*

Proof. Assume that $g_1^m g_2^n = g_2^n g_1^m$. Then the pair $\{g_1, g_2\}$ is not free and the claim follows by Theorem 1 or by Theorem 2. \square

4 Examples

In this section we give examples of commuting morphisms. The morphisms considered in Example 1 can be regarded as direct sums of unary morphisms.

Example 1 Let $X = \{x_1, \dots, x_k\}$ be an alphabet having k letters. Let (m_1, \dots, m_k) and (n_1, \dots, n_k) be k -tuples of nonnegative integers. Define the morphisms $g_1, g_2 \in \text{Tri}(X^*)$ by

$$g_1(x_i) = x_i^{m_i} \quad \text{and} \quad g_2(x_i) = x_i^{n_i}$$

for $i = 1, \dots, k$. Then

$$g_1 g_2(x_i) = g_1(x_i^{n_i}) = x_i^{m_i n_i}$$

and

$$g_2 g_1(x_i) = g_2(x_i^{m_i}) = x_i^{m_i n_i}$$

for $i = 1, \dots, k$. Hence

$$g_1 g_2 = g_2 g_1.$$

Example 2 Let $X = \{a, b\}$ and define the morphisms $g_1, g_2 \in \text{Tri}(X^*)$ by

$$g_1(a) = a, \quad g_1(b) = b^2$$

and

$$g_2(a) = a^2, \quad g_2(b) = b.$$

By Example 1 the morphisms g_1 and g_2 commute. However, there do not exist positive integers m, n and a morphism $g \in \text{Hom}(X^*)$ such that $g_1 = g^m$ and $g_2 = g^n$. To see this, assume that such g, m and n exist. Then neither $g(a)$ nor $g(b)$ is the empty word. Furthermore, either $|g(a)| = 1$ or $|g(b)| = 1$ but not both. Without loss of generality assume that $|g(a)| = 1$. Then $g(a) = a$ or $g(a) = b$. The first alternative is not possible since $g^n(a) = a^2$. The second alternative is not possible since it would imply that the only word of length one in $g(X^*)$ is b .

Example 3 Let $X = \{a, b\}$ and let p and q be positive integers. Let α be a nonnegative integer. Define the morphisms $g_1, g_2 \in \text{Tri}(X^*)$ by $g_1(a) = g_2(a) = a$, $g_1(b) = (ba^\alpha)^{p-1}b$, $g_2(b) = (ba^\alpha)^{q-1}b$.

To prove the equation $g_1 g_2 = g_2 g_1$, let $z = ba^\alpha$. Then $g_1(z) = z^p$ and $g_2(z) = z^q$. Hence $g_1 g_2(z) = g_2 g_1(z)$. Therefore $g_1 g_2(b) a^\alpha = g_2 g_1(b) a^\alpha$, which implies that $g_1 g_2(b) = g_2 g_1(b)$. Trivially $g_1 g_2(a) = g_2 g_1(a)$.

Example 4 Let $X = \{a, b\}$ and let $g_1, g_2 \in \text{Tri}(X^*)$ be nonsingular upper triangular binary morphisms. Assume that there exist positive integers m and n such that $g_1^m = g_2^n$. Assume $g_1 \neq g_2$. Then $\{g_1, g_2\}$ is not a special pair. Indeed, if one of $g_1(a)$ and $g_2(a)$ equals a , then both do. Hence Theorem 1 implies that $g_1 g_2 = g_2 g_1$.

Let u and v be words over the binary alphabet $X = \{a, b\}$. We say that u and v are *a-conjugates* if there exist nonnegative integers p, q, r, s and a word w such that

$$u = a^p w a^q, \quad v = a^r w a^s \quad \text{and} \quad p + q = r + s.$$

Example 5 Let $X = \{a, b\}$ and let $g_1, g_2 \in \text{Tri}(X^*)$ be nonsingular upper triangular morphisms. Assume that $g_1(a) = g_2(a) = a$. Assume that there are positive integers m and n such that $g_1^n(b)$ and $g_2^m(b)$ are *a-conjugates*. We show that these conditions imply that g_1 and g_2 commute. By Lemma 3 it is enough to show that g_1^n and g_2^m commute.

By assumption, there exist nonnegative integers $\gamma_1, \gamma_2, \delta_1, \delta_2, \alpha_1, \dots, \alpha_{p-1}$ such that $g_1^n(b) = a^{\gamma_1} z a^{\gamma_2}$ and $g_2^m(b) = a^{\delta_1} z a^{\delta_2}$, where $z = b a^{\alpha_1} b a^{\alpha_2} b \dots b a^{\alpha_{p-1}} b$ and $\gamma_1 + \gamma_2 = \delta_1 + \delta_2$. Then

$$\begin{aligned} g_1^n g_2^m(b) &= a^{\delta_1} g_1^n(z) a^{\delta_2} \\ &= a^{\delta_1 + \gamma_1} z a^{\gamma_2 + \alpha_1 + \gamma_1} z a^{\gamma_2 + \alpha_2 + \gamma_1} z a^{\gamma_2} \dots a^{\gamma_1} z a^{\gamma_2 + \alpha_{p-1} + \gamma_1} z a^{\gamma_2 + \delta_2} \end{aligned}$$

and

$$\begin{aligned} g_2^m g_1^n(b) &= a^{\gamma_1} g_2^m(z) a^{\gamma_2} \\ &= a^{\gamma_1 + \delta_1} z a^{\delta_2 + \alpha_1 + \delta_1} z a^{\delta_2 + \alpha_2 + \delta_1} z a^{\delta_2} \dots a^{\delta_1} z a^{\delta_2 + \alpha_{p-1} + \delta_1} z a^{\delta_2 + \gamma_2}. \end{aligned}$$

Therefore $g_1^n g_2^m(b) = g_2^m g_1^n(b)$. Hence $g_1^n g_2^m = g_2^m g_1^n$.

We conclude this section by two examples involving singular morphisms.

Example 6 Let $X = \{a, b\}$. Define the morphisms $g_1, g_2 \in \text{Tri}(X^*)$ by

$$g_1(a) = g_2(a) = \varepsilon, \quad g_1(b) = w^i, \quad g_2(b) = w^j$$

where $w \in X^*$ and i and j are nonnegative integers. Then

$$g_1 g_2(b) = g_1(w^j) = w^{ij|w|_b} \quad \text{and} \quad g_2 g_1(b) = g_2(w^i) = w^{ij|w|_b}.$$

Hence $g_1 g_2 = g_2 g_1$.

Example 7 Let $X = \{a, b\}$. Define the morphisms $g_1, g_2 \in \text{Tri}(X^*)$ by

$$g_1(a) = \varepsilon, \quad g_1(b) = (a^\alpha b a^\beta)^i$$

and

$$g_2(a) = a, \quad g_2(b) = (b a^{\alpha+\beta})^j b$$

where α, β, i, j are nonnegative integers. Then

$$g_1 g_2(b) = g_1((b a^{\alpha+\beta})^j b) = g_1(b^{j+1}) = (a^\alpha b a^\beta)^{i(j+1)}$$

and

$$g_2 g_1(b) = g_2((a^\alpha b a^\beta)^i) = (a^\alpha (b a^{\alpha+\beta})^j b a^\beta)^i = (a^\alpha b a^\beta)^{i(j+1)}.$$

Hence $g_1 g_2 = g_2 g_1$.

5 Properties of infinite words generated by upper triangular binary morphisms

Let $X = \{a, b\}$ be a binary alphabet. Regard a as the first letter of X and b as the second letter of X .

Let $h \in \text{Tri}(X^*)$. Assume that h is nonsingular. Then there exist a non-negative integer γ and a word v such that $h(b) = a^\gamma bv$. Let c be a new letter and let $Y = X \cup \{c\}$. Regard c as the third letter of Y . Define the morphism $\mathbf{RIGHT}(h) \in \text{Tri}(Y^*)$ by

$$\mathbf{RIGHT}(h)(x) = h(x), \quad \text{if } x \in X, \quad \mathbf{RIGHT}(h)(c) = cv.$$

Assume that $v \neq \varepsilon$. Then we define the infinite word $\omega(h)$ by

$$\omega(h) = bc^{-1}\mathbf{RIGHT}(h)^\omega(c).$$

In other words, the infinite word $\omega(h)$ is obtained from $\mathbf{RIGHT}(h)^\omega(c)$ by replacing its first letter c by b . Hence, if n is any positive integer, the word obtained from $h^n(b)$ by deleting all occurrences of a preceding the first occurrence of b is a prefix of $\omega(h)$.

For the proof of the following lemma see [6].

Lemma 4 *Let $g_1, g_2 \in \text{Tri}(X^*)$ be nonsingular morphisms. Let $h_i = \mathbf{RIGHT}(g_i)$ for $i = 1, 2$. Assume that $h_i(c) \neq c$ for $i = 1, 2$. If $g_1g_2 = g_2g_1$, then $\omega(g_1) = \omega(g_2)$.*

We will now study some properties of the infinite words defined above.

Let w be an infinite word over X having infinitely many occurrences of b . For $i \geq 1$, let $A_w(i)$ be the number of occurrences of the letter a in w between the i th and the $(i+1)$ th occurrences of b in w .

The following lemma gives a formula for $A_w(i)$, which will be used repeatedly.

Lemma 5 *Let $h \in \text{Tri}(X^*)$ be the morphism defined by*

$$h(a) = a^s \quad \text{and} \quad h(b) = a^{\gamma_1}ba^{\alpha_1}ba^{\alpha_2}b \cdots ba^{\alpha_{p-1}}ba^{\gamma_2},$$

where $s \geq 1$, $p \geq 2$ and $\gamma_1, \gamma_2, \alpha_1, \dots, \alpha_{p-1} \geq 0$. Let $w = \omega(h)$. Then

- (i) $A_w(i + pn) = \alpha_i$ if $i \in \{1, \dots, p-1\}$ and $n \geq 0$.
- (ii) $A_w(pi) = sA_w(i) + \gamma_1 + \gamma_2$ if $i \geq 1$.
- (iii) If $m \geq 1$, $k \geq 0$, $d_m, \dots, d_{m+k} \in \{0, 1, \dots, p-1\}$ and $d_m \neq 0$, then

$$A_w(d_m p^m + d_{m+1} p^{m+1} + \cdots + d_{m+k} p^{m+k}) = \alpha_{d_m} s^m + (\gamma_1 + \gamma_2)(1 + s + \cdots + s^{m-1}).$$

Proof. The infinite word w belongs to $a^{-\gamma_1}h(b)\{a, h(b)\}^\omega$ and $|h(b)|_b = p$. This implies (i).

To prove (ii), let

$$w = w_1 b a^j b \cdots,$$

where $|w_1 b|_b = i$ and $j = A_w(i)$. Then

$$w = a^{-\gamma_1} h(w_1) h(b) a^{js} h(b) \cdots,$$

where $|h(w_1)h(b)|_b = p|w_1 b|_b = pi$. Hence

$$A_w(pi) = \gamma_2 + js + \gamma_1 = sA_w(i) + \gamma_1 + \gamma_2.$$

This proves (ii).

If $m = 1$, (iii) is a consequence of (i) and (ii). Assume inductively that (iii) holds for $m \geq 1$. Assume that $k \geq 0$, $e_{m+1}, \dots, e_{m+1+k} \in \{0, 1, \dots, p-1\}$ and $e_{m+1} \neq 0$. Then

$$\begin{aligned} & A_w(e_{m+1}p^{m+1} + e_{m+2}p^{m+2} + \cdots + e_{m+k+1}p^{m+k+1}) \\ &= sA_w(e_{m+1}p^m + e_{m+2}p^{m+1} + \cdots + e_{m+k+1}p^{m+k}) + \gamma_1 + \gamma_2 \\ &= s(\alpha_{e_{m+1}}s^m + (\gamma_1 + \gamma_2)(1 + s + \cdots + s^{m-1})) + \gamma_1 + \gamma_2 \\ &= \alpha_{e_{m+1}}s^{m+1} + (\gamma_1 + \gamma_2)(1 + s + \cdots + s^m). \end{aligned}$$

Here the first equation follows by (ii) and the second equation by the inductive hypothesis. This proves (iii). \square

The final lemma of this section studies the case of eventually periodic words.

Lemma 6 *Let h be as in Lemma 5. Assume that $w = \omega(h)$ is eventually periodic. Then $\gamma_1 = \gamma_2 = 0$ and $\alpha_1 = \alpha_2 = \cdots = \alpha_{p-1}$.*

Proof. Since w is eventually periodic, also the sequence $(A_w(i))_{i \geq 1}$ is eventually periodic. In particular, this sequence takes only finitely many different values. Hence, by Lemma 5, we have $\gamma_1 = \gamma_2 = 0$. If $\alpha_1 = \cdots = \alpha_{p-1} = 0$, the claim of the lemma holds. Assume that some α_i is nonzero. Then Lemma 5 implies that $s = 1$.

Assume

$$A_w(i) = A_w(i + d) \quad \text{for all } i \geq i_0,$$

where i_0 is an integer and $d = ep^m + fp^{m+1}$ for some $m \geq 0$, $e \in \{1, \dots, p-1\}$ and $f \geq 0$. Choose an integer n such that $n > m$ and $p^n \geq i_0$. Then

$$A_w(jp^n) = A_w(ep^m + fp^{m+1} + jp^n)$$

for $j = 1, \dots, p-1$. Now Lemma 5 implies that

$$\alpha_j = \alpha_e$$

for $j = 1, \dots, p-1$. This implies the claim. \square

6 Commuting nonsingular morphisms g_1 and g_2 such that both $g_1(b)$ and $g_2(b)$ have at least two occurrences of b

In this section $X = \{a, b\}$. We will consider nonsingular morphisms $g_1, g_2 \in \text{Tri}(X^*)$ such that $|g_i(b)|_b \geq 2$ for $i = 1, 2$. We have two different cases to consider according to whether $|g_1(b)|_b$ and $|g_2(b)|_b$ are multiplicatively independent or not. Recall that two integers $p \geq 2$ and $q \geq 2$ are *multiplicatively dependent* if there are positive integers r, m, n such that $p = r^m$ and $q = r^n$ (see [1]).

6.1 The numbers of occurrences of b are multiplicatively independent

Lemma 7 *Let $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, be morphisms such that*

$$g_1(a) = a^s \quad \text{and} \quad g_1(b) = a^{\gamma_1} b a^{\alpha_1} b a^{\alpha_2} b \dots b a^{\alpha_{p-1}} b a^{\gamma_2}$$

and

$$g_2(a) = a^t \quad \text{and} \quad g_2(b) = a^{\delta_1} b a^{\beta_1} b a^{\beta_2} b \dots b a^{\beta_{q-1}} b a^{\delta_2}$$

where $s, t \geq 1$, $p, q \geq 2$ and $\gamma_1, \gamma_2, \delta_1, \delta_2, \alpha_1, \dots, \alpha_{p-1}, \beta_1, \dots, \beta_{q-1} \geq 0$. Assume that p and q are multiplicatively independent. Assume that $g_1(b) \notin b^*$. Assume that $\omega(g_1) = \omega(g_2)$. Then $s = t = 1$ and $\gamma_1 = \gamma_2 = \delta_1 = \delta_2 = 0$.

Proof. Let z be the smallest positive integer such that $\beta_z = \max\{\beta_i \mid i = 1, 2, \dots, q-1\}$. Then $\beta_z \geq 0$ but it is possible that $\beta_z = 0$.

By Lemma 5 we have

$$A_{\omega(g_2)}(i) \leq A_{\omega(g_2)}(zq^n) \tag{1}$$

for $n \geq 1$ and $i < zq^n$. Consider the numbers zq^n , $n \geq 1$. For $n \geq 1$, let

$$zq^n = p^{\tau(n)}(i_n + pj_n),$$

where $\tau(n), j_n \geq 0$ and $i_n \in \{1, \dots, p-1\}$.

Now, the set $\{j_n \mid n \geq 1\}$ is infinite. To see this, assume on the contrary that it is finite. Then there are integers m and n such that $i_m + pj_m = i_n + pj_n$ and $m < n$. This implies that

$$\frac{zq^m}{p^{\tau(m)}} = \frac{zq^n}{p^{\tau(n)}}.$$

Hence $p^{\tau(n)-\tau(m)} = q^{n-m}$, which contradicts the assumption. It follows that the set $\{j_n \mid n \geq 1\}$ is infinite. Therefore there is an integer $n \geq 1$ such that

$$zq^n = p^{\tau(n)}(i_n + x_1p + \dots + x_kp^k)$$

where $k \geq 2$ and $x_k \neq 0$.

Next, let y be an integer such that $\alpha_y = \max\{\alpha_i \mid i = 1, \dots, p-1\}$ and consider the numbers

$$K_1 = yp^{\tau(n)+k-1} \quad \text{and} \quad K_2 = zq^n = p^{\tau(n)}(i_n + x_1p + \dots + x_kp^k).$$

Then we have $K_1 < K_2$. Therefore (1) implies that

$$A_{\omega(g_1)}(K_1) = A_{\omega(g_2)}(K_1) \leq A_{\omega(g_2)}(K_2) = A_{\omega(g_1)}(K_2).$$

On the other hand, Lemma 5 implies that

$$A_{\omega(g_1)}(K_1) = \alpha_y s^{\tau(n)+k-1} + (\gamma_1 + \gamma_2)(1 + s + \dots + s^{\tau(n)+k-2})$$

and

$$A_{\omega(g_1)}(K_2) = \alpha_{i_n} s^{\tau(n)} + (\gamma_1 + \gamma_2)(1 + s + \dots + s^{\tau(n)-1}).$$

Since $A_{\omega(g_1)}(K_1) \leq A_{\omega(g_1)}(K_2)$, we have $\gamma_1 = \gamma_2 = 0$. If $\alpha_y = 0$, we would have $g_1(b) \in b^*$ which contradicts our assumption. Hence $\alpha_y \neq 0$ and $s = 1$.

Since $\gamma_1 = \gamma_2 = 0$ and $s = 1$, we have $A_{\omega(g_1)}(i) \in \{\alpha_1, \dots, \alpha_{p-1}\}$ for all $i \geq 1$. Now the equality $\omega(g_1) = \omega(g_2)$ implies that $\delta_1 = \delta_2 = 0$ and $t = 1$. \square

The next theorem gives all nonsingular morphisms $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, such that $g_1g_2 = g_2g_1$ and the numbers $|g_1(b)|_b$ and $|g_2(b)|_b$ are multiplicatively independent integers larger than one. In the proof we use automatic sequences and Cobham's theorem characterizing sequences which are p -automatic and q -automatic for multiplicatively independent integers p and q (see [1]).

Theorem 8 *Let $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, be morphisms such that*

$$g_1(a) = a^s \quad \text{and} \quad g_1(b) = a^{\gamma_1} b a^{\alpha_1} b a^{\alpha_2} b \dots b a^{\alpha_{p-1}} b a^{\gamma_2}$$

and

$$g_2(a) = a^t \quad \text{and} \quad g_2(b) = a^{\delta_1} b a^{\beta_1} b a^{\beta_2} b \dots b a^{\beta_{q-1}} b a^{\delta_2}$$

where $s, t \geq 1$, $p, q \geq 2$ and $\gamma_1, \gamma_2, \delta_1, \delta_2, \alpha_1, \dots, \alpha_{p-1}, \beta_1, \dots, \beta_{q-1} \geq 0$. Assume that p and q are multiplicatively independent. Then $g_1g_2 = g_2g_1$ if and only if at least one of the following conditions holds:

- (i) $g_i(b) \in b^*$ for $i = 1, 2$,
- (ii) $g_1(a) = g_2(a) = a$, $g_1(b) = (ba^\alpha)^{p-1}b$ and $g_2(b) = (ba^\alpha)^{q-1}b$, where $\alpha = \alpha_1$.

Proof. If (i) or (ii) holds, then $g_1g_2 = g_2g_1$ (see Examples 1 and 3).

Assume that $g_1g_2 = g_2g_1$. By Lemma 4 we have $\omega(g_1) = \omega(g_2)$. Hence, if $g_1(b) \in b^*$, also $g_2(b) \in b^*$ and (i) holds. Assume that $g_1(b) \notin b^*$ and $g_2(b) \notin b^*$.

Now Lemma 7 implies that $s = t = 1$ and $\gamma_1 = \gamma_2 = \delta_1 = \delta_2 = 0$. By Lemma 5, the sequence $(A_{\omega(g_1)}(i))_{i \geq 1}$ is p -automatic and the sequence $(A_{\omega(g_2)}(i))_{i \geq 1}$ is q -automatic. Since these sequences are equal and the numbers p and q are multiplicatively independent, $(A_{\omega(g_1)}(i))_{i \geq 1}$ is eventually periodic. Hence $\omega(g_1)$ is eventually periodic. Now Lemma 6 implies that $\alpha_1 = \alpha_2 = \dots = \alpha_{p-1}$. Hence $g_1(b) = (ba^\alpha)^{p-1}b$ where $\alpha = \alpha_1$.

A similar argument shows that $g_2(b) = (ba^\beta)^{q-1}b$, where $\beta = \beta_1$. Since $\omega(g_1) = \omega(g_2)$, we have $\alpha_1 = \beta_1$. Hence (ii) holds. \square

6.2 The numbers of occurrences of b are multiplicatively dependent

In this subsection we first consider the case that $|g_1(b)|_b$ and $|g_2(b)|_b$ are equal.

Lemma 9 *Let $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, be morphisms such that*

$$g_1(a) = a^s \quad \text{and} \quad g_1(b) = a^{\gamma_1} b a^{\alpha_1} b a^{\alpha_2} b \cdots b a^{\alpha_{p-1}} b a^{\gamma_2}$$

and

$$g_2(a) = a^t \quad \text{and} \quad g_2(b) = a^{\delta_1} b a^{\beta_1} b a^{\beta_2} b \cdots b a^{\beta_{q-1}} b a^{\delta_2}$$

where $s, t \geq 1$, $p, q \geq 2$ and $\gamma_1, \gamma_2, \delta_1, \delta_2, \alpha_1, \dots, \alpha_{p-1}, \beta_1, \dots, \beta_{q-1} \geq 0$. Assume that $p = q$. If $g_1 g_2 = g_2 g_1$, then at least one of the following conditions holds:

- (i) $g_1 = g_2$,
- (ii) $g_i(b) \in b^*$ for $i = 1, 2$,
- (iii) $g_1(a) = g_2(a) = a$ and the words $g_1(b)$ and $g_2(b)$ are a -conjugates.

Proof. Assume that $g_1 g_2 = g_2 g_1$. Let $w_i = \omega(g_i)$ for $i = 1, 2$. By Lemma 4 we have $w_1 = w_2$. Since $p = q$, the equation $w_1 = w_2$ implies that $\alpha_i = \beta_i$ for $i = 1, \dots, p-1$.

If $g_1(b) \in b^*$, we have $w_1 = b^\omega$. This implies that $g_2(b) \in b^*$ and hence (ii) holds.

Assume that $g_1(b) \notin b^*$ and $g_2(b) \notin b^*$.

Next, assume that $A_{w_1}(i)$ takes only finitely many different values. Then Lemma 5 implies that $\gamma_1 = \gamma_2 = 0$. Since $g_1(b) \notin b^*$, some α_i is nonzero. This implies that $s = 1$. By a similar reasoning it is seen that $\delta_1 = \delta_2 = 0$ and $t = 1$. Then $g_1(a) = g_2(a)$ and $g_1(b) = g_2(b)$ and condition (i) holds.

Assume then that $A_{w_1}(i)$ takes infinitely many values. Since $A_{w_1}(pi) = A_{w_2}(pi)$ for all $i \geq 1$, Lemma 5 implies that

$$sA_{w_1}(i) + \gamma_1 + \gamma_2 = tA_{w_2}(i) + \delta_1 + \delta_2$$

for all $i \geq 1$. Because this equation holds for infinitely many different values of $A_{w_1}(i) = A_{w_2}(i)$, it follows that $s = t$ and $\gamma_1 + \gamma_2 = \delta_1 + \delta_2$. If now $s = t = 1$, we have (iii).

Assume that $s = t > 1$. By counting the number of occurrences of a before the first occurrence of b in $g_1 g_2(b) = g_2 g_1(b)$, we see that

$$s\delta_1 + \gamma_1 = t\gamma_1 + \delta_1.$$

By counting the number of occurrences of a after the last occurrence of b in $g_1 g_2(b) = g_2 g_1(b)$, we see that

$$\gamma_2 + s\delta_2 = \delta_2 + t\gamma_2.$$

Hence $(s-1)\delta_1 = (t-1)\gamma_1$ and $(s-1)\delta_2 = (t-1)\gamma_2$. Since $s = t > 1$ we have $\gamma_1 = \delta_1$ and $\gamma_2 = \delta_2$. Therefore condition (i) holds. \square

The next theorem gives all nonsingular morphisms $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, such that $g_1 g_2 = g_2 g_1$ and the numbers $|g_1(b)|_b$ and $|g_2(b)|_b$ are multiplicatively dependent integers larger than one.

Theorem 10 Let $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, be morphisms such that

$$g_1(a) = a^s \quad \text{and} \quad g_1(b) = a^{\gamma_1} b a^{\alpha_1} b a^{\alpha_2} b \cdots b a^{\alpha_{p-1}} b a^{\gamma_2}$$

and

$$g_2(a) = a^t \quad \text{and} \quad g_2(b) = a^{\delta_1} b a^{\beta_1} b a^{\beta_2} b \cdots b a^{\beta_{q-1}} b a^{\delta_2}$$

where $s, t \geq 1$, $p, q \geq 2$ and $\gamma_1, \gamma_2, \delta_1, \delta_2, \alpha_1, \dots, \alpha_{p-1}, \beta_1, \dots, \beta_{q-1} \geq 0$. Assume that $p = r^m$ and $q = r^n$ where m, n, r are positive integers. Then $g_1 g_2 = g_2 g_1$ if and only if at least one of the following conditions holds:

- (i) $g_1^n = g_2^m$,
- (ii) $g_i(b) \in b^*$ for $i = 1, 2$,
- (iii) $g_1(a) = g_2(a) = a$ and the words $g_1^n(b)$ and $g_2^m(b)$ are a -conjugates.

Proof. If at least one of the conditions (i), (ii) or (iii) holds, then $g_1 g_2 = g_2 g_1$ (see Examples 1, 4 and 5).

Assume $g_1 g_2 = g_2 g_1$. Let $h_1 = g_1^n$ and $h_2 = g_2^m$. Then $|h_1(a)|_a = s^n$, $|h_2(a)|_a = t^m$ and $|h_1(b)|_b = p^n = r^{mn} = q^m = |h_2(b)|_b$.

Since $h_1 h_2 = h_2 h_1$, Lemma 9 implies that at least one of the following conditions holds:

- (a) $h_1 = h_2$,
 - (b) $h_i(b) \in b^*$ for $i = 1, 2$,
 - (c) $h_1(a) = h_2(a) = a$ and the words $h_1(b)$ and $h_2(b)$ are a -conjugates.
- Now (a) implies (i), (b) implies (ii) and (c) implies (iii). \square

7 Commuting nonsingular morphisms g_1 and g_2 such that $|g_1(b)|_b = 1$ and $|g_2(b)|_b \geq 2$

Let $X = \{a, b\}$. The following theorem gives all commuting nonsingular morphisms $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, such that $|g_1(b)|_b = 1$ and $|g_2(b)|_b \geq 2$.

Theorem 11 Let $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, be morphisms such that

$$g_1(a) = a^s \quad \text{and} \quad g_1(b) = a^{\gamma_1} b a^{\gamma_2}$$

and

$$g_2(a) = a^t \quad \text{and} \quad g_2(b) = a^{\delta_1} b a^{\beta_1} b a^{\beta_2} b \cdots b a^{\beta_{q-1}} b a^{\delta_2}$$

where $s, t \geq 1$, $q \geq 2$ and $\gamma_1, \gamma_2, \delta_1, \delta_2, \beta_1, \dots, \beta_{q-1} \geq 0$. Then $g_1 g_2 = g_2 g_1$ if and only if at least one of the following conditions holds:

- (i) $g_1(x) = x$ for all $x \in \{a, b\}$,
- (ii) $g_i(b) \in b^*$ for $i = 1, 2$.

Proof. If (i) or (ii) holds, then $g_1 g_2 = g_2 g_1$.

Assume then that $g_1 g_2 = g_2 g_1$. Let $h_1 = g_1 g_2$ and $h_2 = g_2$. Then $h_1 h_2 = h_2 h_1$. Since $|h_1(b)|_b = |h_2(b)|_b \geq 2$, Lemma 9 implies that at least one of the following conditions holds:

- (a) $h_1 = h_2$,
 - (b) $h_i(b) \in b^*$ for $i = 1, 2$,
 - (c) $h_1(a) = h_2(a) = a$ and the words $h_1(b)$ and $h_2(b)$ are a -conjugates.
- Now (a) implies (i), (b) implies (ii) and (c) implies (i). \square

8 Commuting nonsingular morphisms g_1 and g_2 such that $|g_1(b)|_b = |g_2(b)|_b = 1$

Let $X = \{a, b\}$. In this section we give all commuting nonsingular morphisms $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, such that $|g_1(b)|_b = |g_2(b)|_b = 1$.

Proposition 12 *Let $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, be morphisms such that*

$$g_1(a) = a^s \quad \text{and} \quad g_1(b) = a^{\gamma_1} b a^{\gamma_2}$$

and

$$g_2(a) = a^t \quad \text{and} \quad g_2(b) = a^{\delta_1} b a^{\delta_2}$$

where $s, t \geq 1$ and $\gamma_1, \gamma_2, \delta_1, \delta_2 \geq 0$. Then $g_1 g_2 = g_2 g_1$ if and only if

$$(s-1)\delta_i = (t-1)\gamma_i \quad \text{for} \quad i = 1, 2.$$

Proof. We have

$$g_1 g_2(b) = a^{s\delta_1 + \gamma_1} b a^{s\delta_2 + \gamma_2} \quad \text{and} \quad g_2 g_1(b) = a^{t\gamma_1 + \delta_1} b a^{t\gamma_2 + \delta_2}.$$

Hence $g_1 g_2 = g_2 g_1$ if and only if $s\delta_1 + \gamma_1 = t\gamma_1 + \delta_1$ and $s\delta_2 + \gamma_2 = t\gamma_2 + \delta_2$. This implies the claim. \square

9 Commuting morphisms g_1 and g_2 such that g_1 or g_2 is singular

Let $X = \{a, b\}$ and let $h \in \text{Tri}(X^*)$. If h is singular, then $h(a) = \varepsilon$ or $h(b) \in a^*$. In this section we give all commuting morphisms $g_i \in \text{Tri}(X^*)$, $i = 1, 2$, such that g_1 or g_2 is singular.

Proposition 13 *Let $g_i \in \text{Tri}(X^*)$ for $i = 1, 2$. Assume that $g_1(b) \in a^*$. Then $g_1 g_2 = g_2 g_1$ if and only if $|g_1 g_2(b)| = |g_2 g_1(b)|$ or, equivalently,*

$$|g_1(a)||g_2(b)|_a + |g_1(b)||g_2(b)|_b = |g_2(a)||g_1(b)|.$$

Proof. Since $g_2 g_1(b) \in a^*$ and $g_1 g_2(b) \in a^*$, the claim holds. \square

Proposition 14 *Let $g_1, g_2 \in \text{Tri}(X^*)$ be morphisms such that*

$$g_1(a) = \varepsilon, \quad g_1(b) = u$$

and

$$g_2(a) = a^t, \quad g_2(b) = v,$$

where $t \geq 0$ and both u and v have at least one occurrence of b . Then $g_1g_2 = g_2g_1$ if and only if at least one of the following conditions holds:

- (i) $g_1 = g_2$,
- (ii) $g_2(x) = x$ for all $x \in X$,
- (iii) $t = 0$ and $uv = vu$,
- (iv) $g_i(b) \in b^*$ for $i = 1, 2$,
- (v) $t = 1$ and there exist nonnegative integers α, β, i and j such that

$$g_1(b) = (a^\alpha b a^\beta)^i, \quad g_2(b) = (b a^{\alpha+\beta})^j b.$$

Proof. If (i), (ii), (iii), (iv) or (v) holds, then $g_1g_2 = g_2g_1$ (see Examples 1, 6, 7).

Assume that $g_1g_2 = g_2g_1$. Then

$$g_2(u) = g_2g_1(b) = g_1g_2(b) = g_1(v) = u^{|v|_b}.$$

If $u \in b^*$, this equation implies that $v \in b^*$. Hence (iv) holds. Assume that $u \notin b^*$.

Next, assume that $t = 0$. Then $u^{|v|_b} = g_2(u) = v^{|u|_b}$, which shows that (iii) holds.

Assume then that $t \neq 0$. By assumption, $|v|_b = 1$ or $|v|_b \geq 2$. Assume first that $|v|_b = 1$. Then the equation $g_2(u) = u$ shows that (ii) holds. Assume finally that $|v|_b \geq 2$. Then $\omega(g_2)$ is defined. Since $\omega(g_2)$ is obtained from $g_2^\omega(u) = u^\omega$ by deleting the occurrences of a preceding the first occurrence of b , the word $\omega(g_2)$ is eventually periodic. Hence $t = 1$. By Lemma 6 there are nonnegative integers γ and j such that $g_2(b) = (b a^\gamma)^j b$. Since $\omega(g_2) = (b a^\gamma)^\omega$, there exist nonnegative integers α, β and i such that $g_1(b) = (a^\alpha b a^\beta)^i$ and $\gamma = \alpha + \beta$. Hence (v) holds. \square

For a systematic study of the equation $h(w) = w^n$, $n \geq 2$, for binary morphisms, see [3].

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